INDIAN INSTITUTE OF SCIENCE E9-252: MATHEMATICAL METHODS AND TECHNIQUES IN SIGNAL PROCESSING HOME WORK #2 - SOLUTIONS, FALL 2014

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Problem 1. (Frequency domain analysis) From Figure 1,

$$X_{1}(z) = \frac{1}{M} \sum_{i=0}^{M-1} X\left(z^{\frac{1}{M}} e^{j\frac{2\pi i}{M}}\right)$$

$$\implies Y_{1}(z) = X_{1}\left(z^{L}\right)$$

$$= \frac{1}{M} \sum_{i=0}^{M-1} X\left(z^{\frac{L}{M}} e^{j\frac{2\pi i}{M}}\right).$$
(1)

Similarly,

$$X_{2}(z) = X(z^{L})$$

$$Y_{2}(z) = \frac{1}{M} \sum_{i=0}^{M-1} X_{2} \left(z^{\frac{1}{M}} e^{j\frac{2\pi i}{M}} \right)$$

$$= \frac{1}{M} \sum_{i=0}^{M-1} X \left(\left(z^{\frac{1}{M}} e^{j\frac{2\pi i}{M}} \right)^{L} \right)$$

$$= \frac{1}{M} \sum_{i=0}^{M-1} X \left(z^{\frac{L}{M}} e^{j\frac{2\pi iL}{M}} \right).$$
(2)



FIGURE 1. Comparing the outputs by changing the order of decimator and upsampler.

To prove that $Y_1(z) = Y_2(z) \forall X(z)$, it is necessary and sufficient to satisfy the following condition:

$$\left\{ X \left(z^{\frac{L}{M}} e^{j\frac{2\pi i L}{M}} \right) \mid i = 0, 1, \cdots, M - 1 \right\} = \left\{ X \left(z^{\frac{L}{M}} e^{j\frac{2\pi i i}{M}} \right) \mid i = 0, 1, \cdots, M - 1 \right\} \quad \forall X (z)$$

$$i.e., \left\{ e^{j\frac{2\pi i L}{M}} \mid i = 0, 1, \cdots, M - 1 \right\} = \left\{ e^{j\frac{2\pi i i}{M}} \mid i = 0, 1, \cdots, M - 1 \right\}.$$

Since $e^{j2\pi k} = 1 \forall k \in \mathbb{Z}$, we have $e^{j\frac{2\pi iL}{M}} = e^{j\frac{2\pi (iL \mod M)}{M}}$. Hence, the equivalent condition is

$$\{(iL) \mod M \mid i = 0, 1, \cdots, M - 1\} = \{0, 1, \cdots, M - 1\}.$$
(3)

Let $0 \le i_1 \le M - 1$ and $0 \le i_2 \le M - 1$ such that $i_1 \ne i_2$. Without loss of generality, consider $i_1 < i_2$. Using the following identity on modulo operation

$$(a-b) \mod M = (a \mod M - b \mod M) \mod M$$
,

we have,

$$((i_1L) \mod M - (i_2L) \mod M) \mod M = ((i_1 - i_2)L) \mod M.$$
(4)

Case L and M are relatively prime:

Since $0 < i_1 - i_2 < M$, and gcd(L, M) = 1, $((i_1 - i_2)L) \mod M \neq 0$. Therefore from (4),

 $((i_1L) \mod M - (i_2L) \mod M) \mod M \neq 0,$

 \implies $(i_1L) \mod M \neq (i_2L) \mod M$.

We have proved that $i_1 \neq i_2 \implies (i_1L) \mod M \neq (i_2L) \mod M \forall i_1, i_2 \in \{0, 1, 2 \cdots, M-1\}$. Therefore, when gcd(L, M) = 1, equation (3) holds true.

Case M divides L:

Let $L = P \times M$, P > 1. Therefore, it is possible to chose $i_1 = i_2 + M$. Under this condition,

$$((i_1 - i_2)L) \mod M = (ML) \mod M = 0.$$

Therefore,

$$\begin{array}{l} ((i_1L) \mod M - (i_2L) \mod M) \mod M = 0 \\ \implies (i_1L) \mod M = (i_2L) \mod M. \end{array}$$

We have shown that for some choice of $i_1 \neq i_2$, $(i_1L) \mod M = (i_2L) \mod M$. Hence, the values $\{(iL) \mod M\}_{i=0}^{M-1}$ are not distinct. Therefore, when M divides L, equation (3) does not hold true.

Case gcd(M, L) = G > 1:

Let $M = G \times P_M$ and $L = G \times P_L$. We can chose $i_1 = i_2 + G$. Under this condition, $e^{j2\pi \frac{iL}{M}} = e^{j2\pi \frac{iP_L}{P_M}}$. Therefore, $\left\{e^{j2\pi \frac{iP_L}{P_M}} \mid i = 0, 1, \cdots, M - 1\right\}$ has P_M distinct values. Therefore, equation (3) does not hold true under this condition.

Hence, the equation (3) holds true iff L and M are relatively prime. This proves that M fold decimator and L fold upsampler blocks can be interchanged iff L and M are relatively prime.

(Time domain analysis) From the definitions of decimator and upsampler,

$$\begin{aligned} x_1[n] &= x[Mn]. \\ y_1[n] &= \begin{cases} x_1\left[\frac{n}{L}\right], & n \text{ is a multiple of } L \\ 0 & \text{otherwise,} \end{cases} \\ y_1[n] &= \begin{cases} x\left[M\frac{n}{L}\right], & n \text{ is a multiple of } L \\ 0 & \text{otherwise.} \end{cases} \end{aligned}$$
 (5)

Similarly,

$$x_{2}[n] = \begin{cases} x_{1}\left[\frac{n}{L}\right], & n \text{ is a multiple of } L \\ 0 & \text{otherwise.} \end{cases}$$

$$y_{1}[n] = x_{1}[Mn],$$

$$y_{1}[n] = \begin{cases} x\left[\frac{Mn}{L}\right], & Mn \text{ is a multiple of } L \\ 0 & \text{otherwise.} \end{cases}$$
(6)

From equations (5) and (6), the outputs are same iff n is a multiple of L when ever Mn is a multiple of L. Case gcd(L, M) = 1: Trivial in this case that L divides $Mn \iff L$ divides n.

Case $gcd(L, M) = P \neq 1$: Let $L = P \times Q$. In this case L divides Mn when ever Q divides n. Hence L divides $Mn \neq L$ divides n.

Therefore, the outputs are same iff L and M are relatively prime.

Problem 2. We use the identities in Figure 2 to simplify the given transformations.



FIGURE 2. Identities related to decimation, upsampling and delay operations.







Problem 3. Figure 3 indicates the stopband and passband frequencies for the lowpass filter when the signal is bandlimited to ω_B . We have

$$\omega_s = \frac{2\pi - \omega_B}{L} = 0.2742$$
$$\omega_p = \frac{\omega_B}{L} = 0.04.$$



FIGURE 3. The frequency response at the output of upsampler is shown. The stopband and passband frequencies for the filter are also indicated.

The stopband and passband ripple amplitudes are

$$\delta_s = 0.005,$$

$$\delta_p = 0.01.$$

Let N be the filter order of H(z) and f_s indicate the sampling frequency at the output of H(z). The filter does N multiplications to give one output sample. Hence, the computational complexity of the efficient implementation of the interpolation filter is Nf_s multiplications per second.

The filter order N is estimated based on the following empirical formula by Herrmann et al.¹

$$N = \frac{D_{\infty}(\delta_p, \delta_s)}{(\omega_s - \omega_p)/2\pi}$$
$$D_{\infty}(\delta_p, \delta_s) = (\log_{10} \delta_s) \left[a_1 \left(\log_{10} \delta_p \right)^2 + a_2 \log_{10} \delta_p + a_3 \right]$$
$$+ \left[a_4 \left(\log_{10} \delta_p \right)^2 + a_5 \log_{10} \delta_p + a_6 \right]$$

where $a_1 = 5.3e - 3$, $a_2 = 0.071$, $a_3 = -0.4761$, $a_4 = -0.0026$, $a_5 = -0.5941$, $a_6 = -0.4278$. A simpler but less accurate empirical formula is given by Bellanger²:

$$N = \frac{2\log_{10}\left(1/\delta_s\delta_p\right)}{3\left(\omega_s - \omega_p\right)/2\pi}.$$

For a signal bandlimited to 3 KHz, the Nyquist sampling rate is 6 KHz and 20% oversampling is at 7.2 KHz. Hence,

$$f_s = 7200 \text{ Hz}$$



FIGURE 4. 2 stage interpolation filter.

For the 2-stage interpolator shown in Figure 4, for the filter $H_1(z)$, let the stopband and passband frequencies be ω_{s1} and ω_{p1} . Let the stopband and passband ripples be δ_{s1} and δ_{p1} respectively. Let the input to $H_1(z)$ be

¹O. Herrmann et al., "Practical design rules for optimum low pass FIR digital filters", Bell-sys tech. Journal, vol 52, no.2, July 1973.

²M. Bellanger, "On computational complexity in digital filters,"Proc. *The Eurioeab Conference on Circuit Theory & Design*, The Haugue, The Netherlands, pp. 58-63, August 1981.

bandlimited to ω_{B1} and let f_{s1} be the sampling frequency at the output of $H_1(z)$ and N_1 be the filter order. We have

$$\begin{split} \omega_{B1} &= \omega_B \\ \omega_{s1} &= \frac{2\pi - \omega_{B1}}{L_1} \\ \omega_{p1} &= \frac{\omega_{B1}}{L_1} \\ f_{s1} &= f_s L_1. \end{split}$$

Let the $\omega_{s2}, \omega_{p2}, \delta_{s2}, \delta_{p2}, \omega_{B2}, f_{s2}, N_2$ represent the corresponding parameters for the second stage filter $H_2(z)$. Since the output of $H_1(z)$, is bandlimited to $\frac{\omega_{B1}}{L_1}$, we have

$$\omega_{B2} = \frac{\omega_{B1}}{L_1}$$
$$\omega_{s2} = \frac{2\pi - \omega_{B2}}{L_2}$$
$$\omega_{p2} = \frac{\omega_{B2}}{L_1}$$
$$f_{s2} = f_s L_2.$$

For cascaded filters, the passband ripples get added and the stopband ripples remain the same. Hence, we choose

$$\delta_{s1} = \delta_{s2} = \delta_s = 0.005,$$

 $\delta_{p1} = \delta_{p2} = \frac{\delta_p}{2} = 0.005.$

Table 1, gives the parameters for various choices of L_1 and L_2 . Table 2 compares the filters orders and computational complexities for various realizations of the interpolation filter. From the table, the two stage implementation with 5-fold interpolation filter followed by 4-fold interpolation filter gives the best performance. This combination is ≈ 4.1 times faster compared to the 1-stage implementation.

(L_1, L_2)	1-stage	(2, 10)	(10, 2)	(4,5)	(5, 4)
$\omega_{B1} = \omega_B$	0.8	0.8	0.8	0.8	0.8
$\omega_{s1} = \frac{2\pi - \omega_{B1}}{L_1}$	0.2742	2.7416	0.5483	1.3708	1.9066
$\omega_{p1} = \frac{\omega_{B1}}{L_1}$	0.04	0.4	0.08	0.2	0.16
$f_{s1} = f_s \times L_1 \text{ (Hz)}$	144000	14400	72000	28800	36000
$\omega_{B2} = \omega_{B1}$	-	0.4	0.08	0.2	0.16
$\omega_{s2} = \frac{2\pi - \omega_{B2}}{L_2}$	-	0.5883	3.1016	1.2166	1.5308
$\omega_{p2} = \frac{2\pi - \bar{\omega}_{B2}}{L_2}$	-	0.04	0.04	0.04	0.04
$f_{s2} = f_{s1} \times L_2 \text{ (Hz)}$	-	144000	144000	144000	144000

TABLE 1. Filter design parameters for 20 fold 2 stage interpolation filters. L_1 fold interpolation filter is followed by L_2 fold interpolation filter.

	(L_1, L_2)	1-stage	(2, 10)	(10, 2)	(4,5)	(5, 4)
Herrmann et al.	Filter order N_1	57	7	31	13	16
	complexity $C_1 = f_{s1}N_1$ of $H_1(z)$ (mult/sec)	8208000	100800	2304000	50400	576000
	Filter order N_2	-	27	5	13	10
	complexity $C_2 = f_{s2}N_2$ of $H_2(z)$	-	3888000	720000	1872000	1440000
	Overall complexity $C_1+C_2(\text{mult/sec})$	8208000	3988800	3024000	2246400	2016000
Bellanger' formula	Filter order N_1	77	9	42	17	21
	complexity $C_1 = f_{s1}N_1$ of $H_1(z)$ (mult/sec)	11088000	129600	3024000	489600	756000
	Filter order N_2	-	36	7	17	13
	complexity $C_2 = f_{s2}N_2$ of $H_2(z)$	-	5184000	1008000	2448000	1872000
	Overall complexity $C_1+C_2(\text{mult/sec})$	11088000	5313600	4032000	2937600	2628000

TABLE 2. Comparison of filter orders and computational complexity of 20 fold 2 stage interpolation filters. $\delta_{s1} = \delta_{s2} = 0.005$. $\delta_{p1} = \delta_{p2} = 0.005$. $D_{\infty} (\delta_{p1}, \delta_{s1}) = D_{\infty} (\delta_{p2}, \delta_{s2}) = 2.3032$. $D_{\infty} (\delta_p, \delta_s) = 2.135$

Problem 4. For the filter H(z), we identify $H_{0}(z)$, $H_{1}(z)$ and $H_{2}(z)$ such that

$$H(z) = H_0(z^3) + z^{-1}H_1(z^3) + z^{-2}H_1(z^3).$$

We identify $H_{00}(z), H_{01}(z), H_{10}(z), H_{11}(z), H_{20}(z), H_{21}(z)$ such that

$$\begin{array}{rcl} H_0\left(z\right) &=& H_{00}(z^2)+z^{-1}H_{01}(z^2),\\ H_1\left(z\right) &=& H_{10}(z^2)+z^{-1}H_{11}(z^2),\\ H_2\left(z\right) &=& H_{20}(z^2)+z^{-1}H_{21}(z^2). \end{array}$$

The efficient implementation of fractional sampling rate alteration starting from decimation filter is as follows.



Let F be the sample rate of x[n] and N be the order of the filter H(z). Note that the sum of filter orders of $\{H_{ij}(z)\}_{i=0,1,2;j=0,1}$ is equal to N. Also, the filters $H_{ij}(z)$ each operates at the rate of F/3 samples/sec. Hence the overall computational complexity is $N\frac{F}{3}$ multiplications/second. The efficient implementation derived in the class starting from the interpolation stage has the same number of

The efficient implementation derived in the class starting from the interpolation stage has the same number of filters operating at the rate of F/3 samples per second. The overall computational complexity in this case is also $N\frac{F}{3}$ multiplications/second.