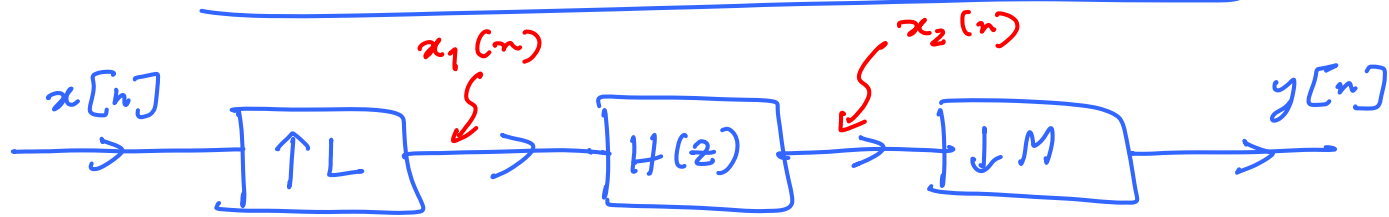


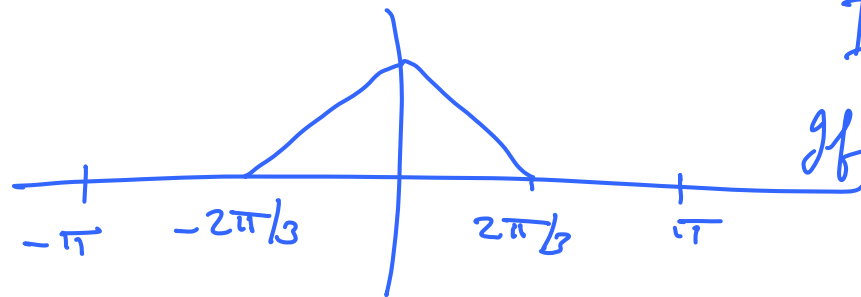
Fractional sampling rate alterations



The effective sampling rate conversion is $\frac{L}{M}$.

Example:

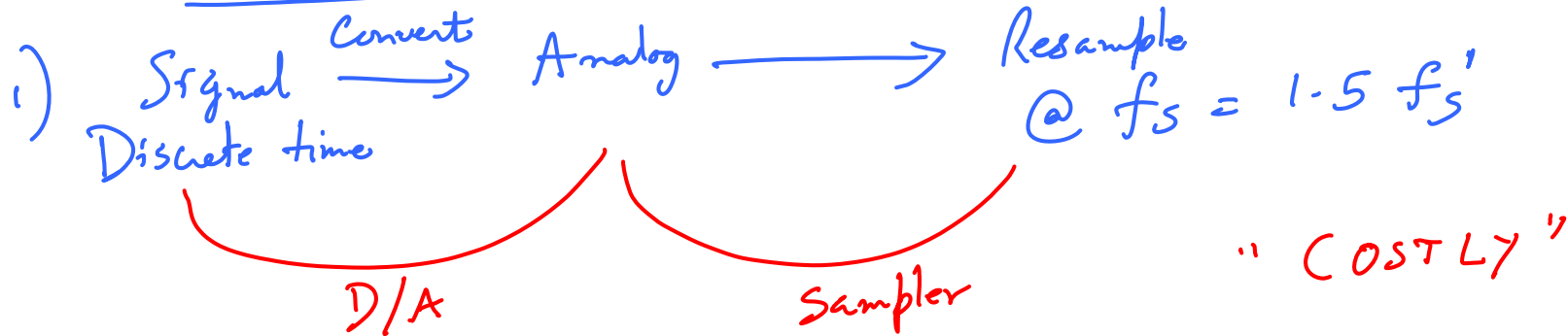
Suppose we have a signal such that $X(e^{j\omega})$ is band limited to $|\omega| < 2\pi/3$



If we decimate by 2 \rightarrow "aliasing"
 If we alter the rate by 1.5 \rightarrow "no aliasing critically"
 $2\pi/3 \cdot 3/2 \rightarrow \pi$

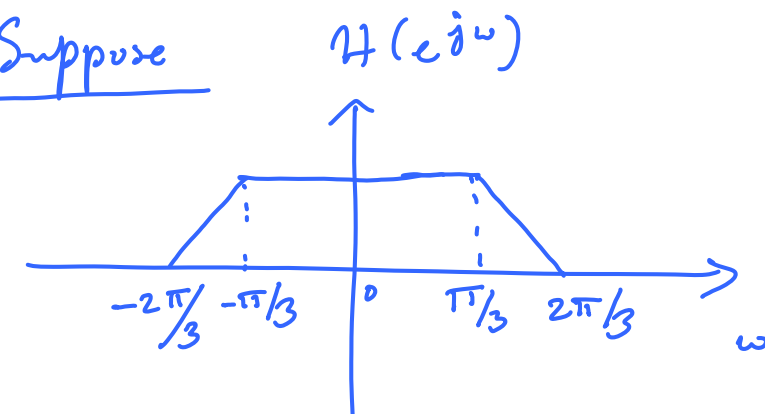
2 approaches for sampling rate conversion

$\frac{L}{M} = 1.5$ say



2) Do a "judicious" use of 'downsampling' & 'upsampling'

Suppose



In general, we can reduce the sampling rate by any rational no M/L .

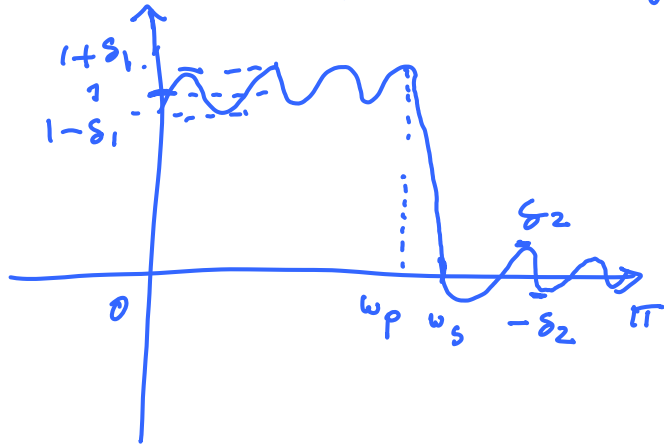
Quality of the filter $H(z)$ determines the quality of the o/p.

Stretched versions of $X(e^{j\omega L/M})$ govern the time domain meaning of L/M . $L > M$ is very much possible

Focussing on the example

$$\begin{aligned} \text{Transition bandwidth} &= \pi/3 \\ \text{normalized tran. bw.} &= \frac{\pi/3}{2\pi} = \frac{1}{6} \\ &(\Delta f) \end{aligned}$$

For an equiripple design,



$$N \approx \frac{2 \log_{10} \left(\frac{1}{10 \delta_1 \delta_2} \right)}{3 \Delta f}$$

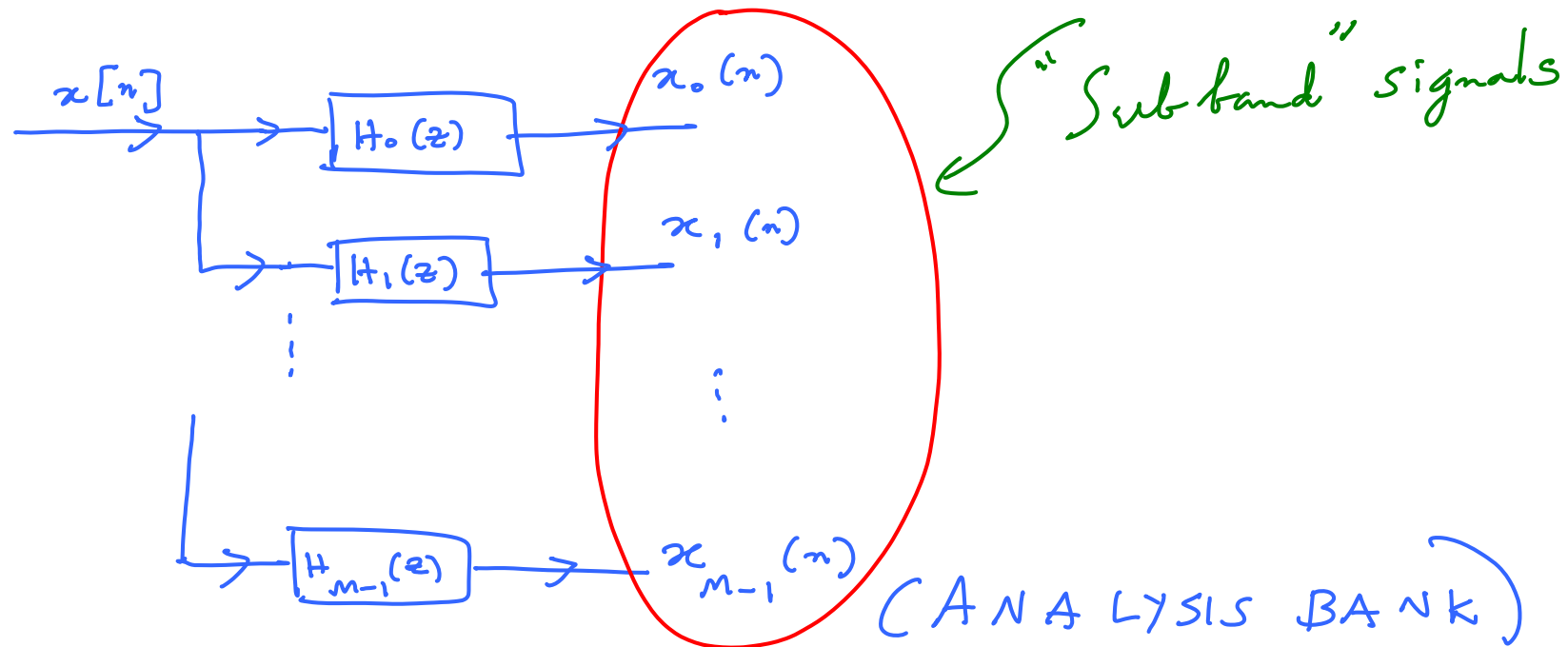
↑
normalized transition
bandwidth

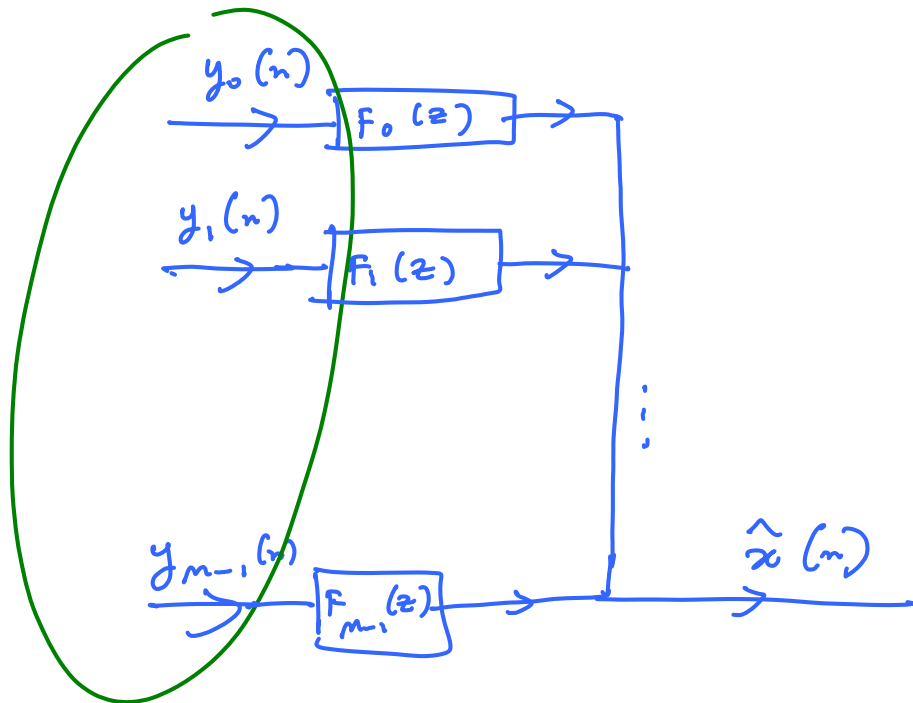
Suppose $\delta_1 = \delta_2 = 0.01$ for
an equiripple filter, $\xi \Delta f = \frac{1}{6}$

$$N \approx 11 \quad (\text{11}^{\text{th}} \text{ order filter})$$

Digital Filter Banks

A digital filter bank is a collection of filters with a common i/p & a common o/p.

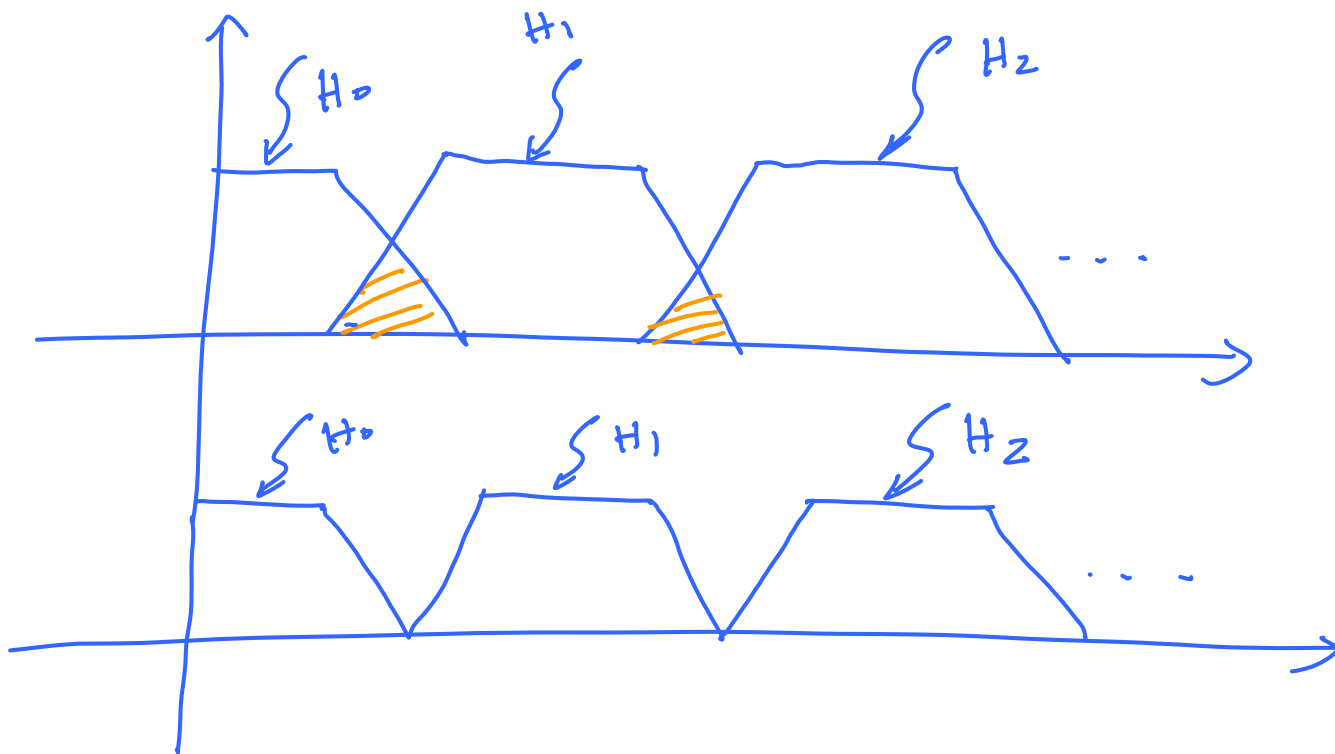




Subband signals

"SYNTHESIS BANK"

What could be the nature of freq. responses for $H_i(z)$



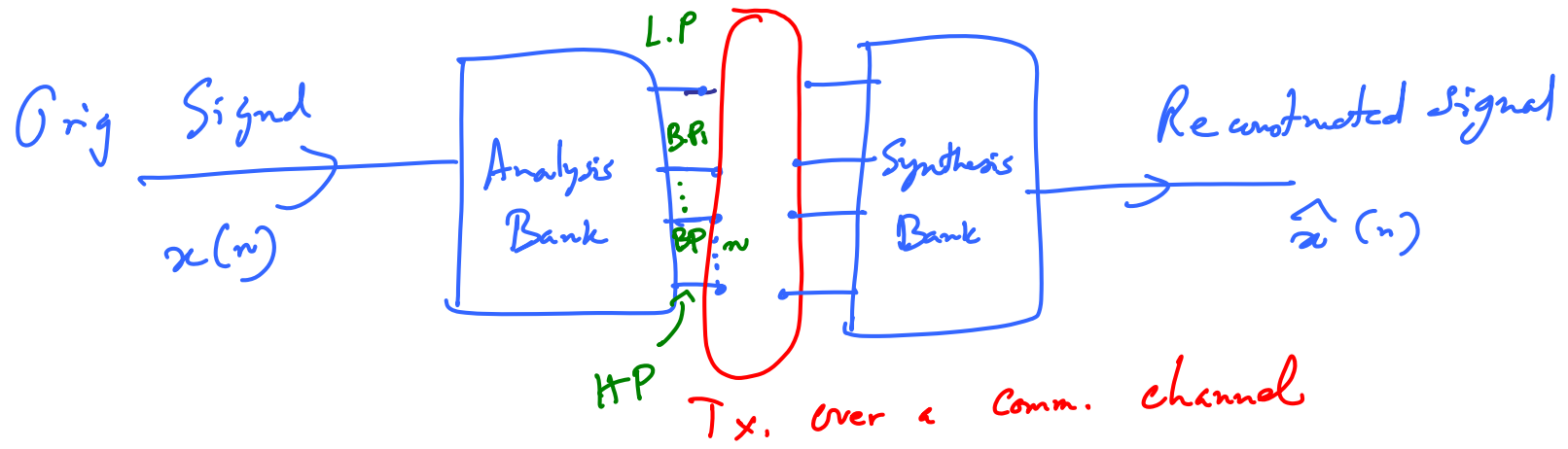
"Overlapping"

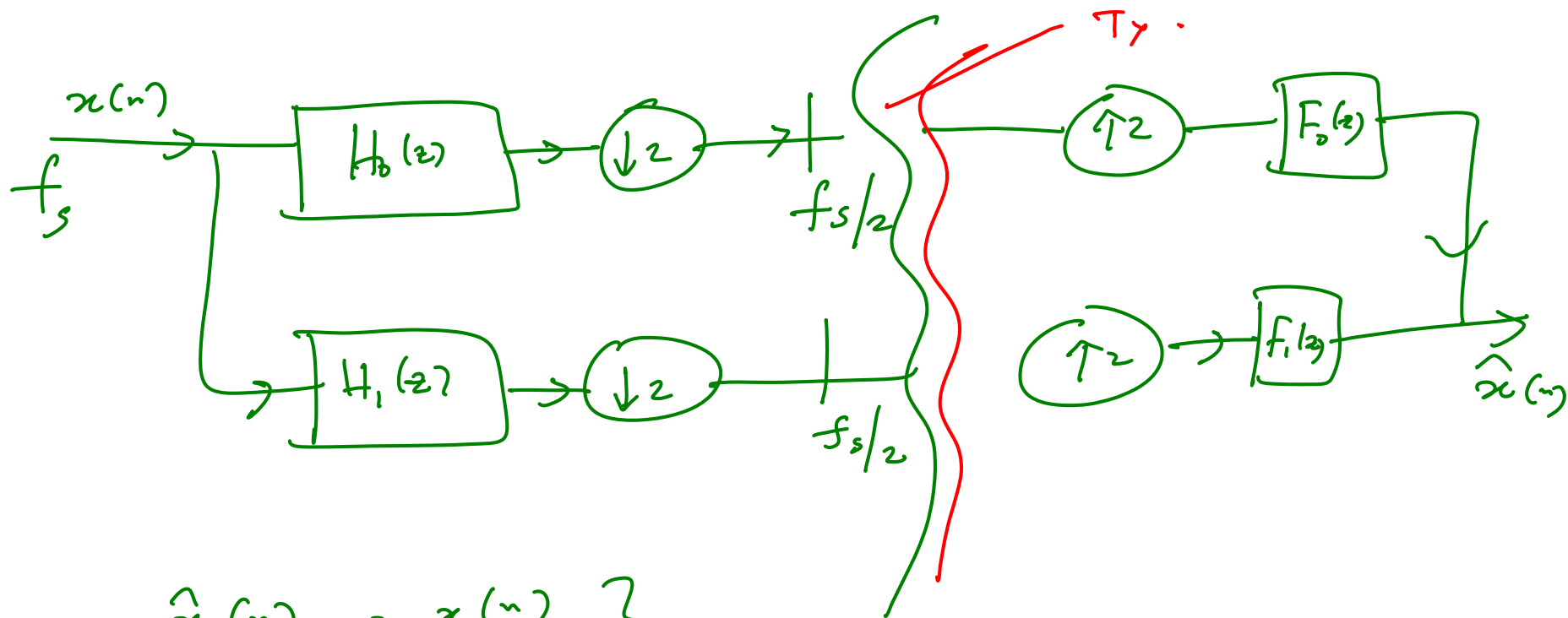
more relaxed slope
in the transition bandwidth
region → Easy to design
filters
Aliasing is an issue

"Non overlapping"

Steeper freq. transition
bandwidth slope

→ Filter orders could be
prohibitive
→ "No aliasing"





Under
what cond

$$\hat{x}(n) = x(n) ?$$

DFT as a simplest filter bank

The DFT matrix ($N \times N$) is such that

$$W_N := \begin{bmatrix} \omega_N^{km} \end{bmatrix} \quad \omega_N = e^{-j2\pi/N}$$

$$X(k) = \sum_{m=0}^{N-1} x(m) \omega_N^{km} \quad (\text{DFT})$$

$$x(m) = \frac{1}{N} \sum_{k=0}^{N-1} X(k) \omega_N^{-km} \quad (\text{IDFT})$$

The entry in the k^{th} row & m^{th} column is $e^{-j \frac{2\pi km}{N}}$

$$\underbrace{X}_{\substack{\text{freq. domain} \\ \text{samples}}} = \underbrace{W_N}_{\text{DFT matrix}} \underbrace{x}_{\substack{\text{time domain} \\ \text{samples}}}$$

ignore this subscript if you wish

Example:

$$N = 2$$

$$W_2 =$$

$$\begin{bmatrix} e^{-j\frac{2\pi}{2} \cdot 0 \cdot 0} & 0 \\ e^{-j\frac{2\pi}{2} \cdot 1 \cdot 0} & 0 \end{bmatrix}$$

$$\begin{bmatrix} e^{-j\frac{2\pi}{2} \cdot 0 \cdot 1} & 0 \\ e^{-j\frac{2\pi}{2} \cdot 1 \cdot 1} & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

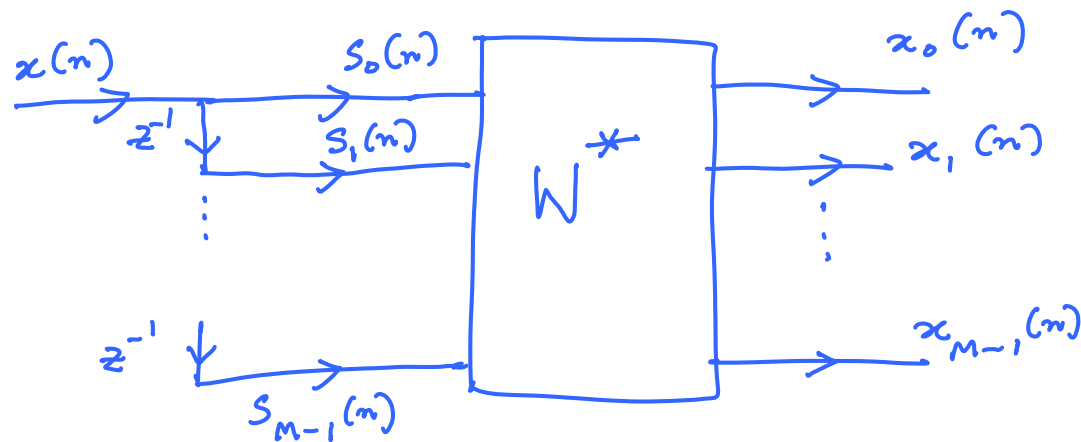
low pass

$$1, 1 \rightarrow 1 + z^{-1} \text{ (LP)}$$

$$1, -1 \rightarrow 1 - z^{-1} \text{ (HP)}$$

From the definitions of W , we can interpret
DFT as a filter bank

Consider the sequence $x(n)$ from which we generate M sequences $s_i(n)$; $i = 0, 1, \dots, M-1$ by passing $x(n)$ through a delay line such that $s_i(n) = x(n-i)$



$$x_k(n) = \sum_{i=0}^{M-1} s_i(n) w_M^{-ki}$$

This is like the IDFT without a scale factor of $\frac{1}{M}$.

Taking Z-transforms,

$$X_k(z) = \sum_{i=0}^{M-1} S_i(z) w_M^{-ki}$$

But $S_i(z) = z^{-i} X(z)$

$$X_k(z) = \sum_{i=0}^{M-1} z^{-i} w_M^{-ki} X(z)$$

$$= \sum_{i=0}^{M-1} \left(z w_M^k \right)^{-i} X(z)$$

$$X_k(z) = H_k(z) X(z)$$

$$H_k(z) = H_0(z w_M^k) ; H_0(z) = 1 + z^{-1} + z^{-2} + \dots + z^{-(M-1)}$$

← modulation term
← mother filter

The DFT system is like a filter bank with analysis filters $H_k(z)$

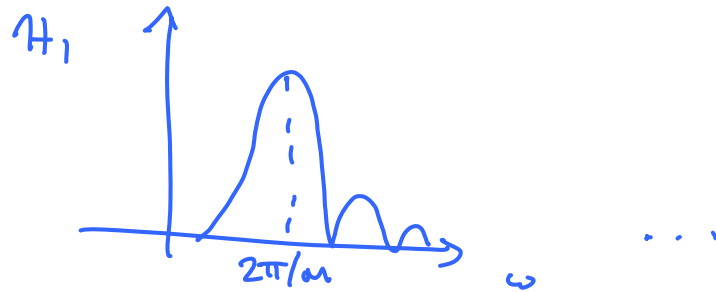
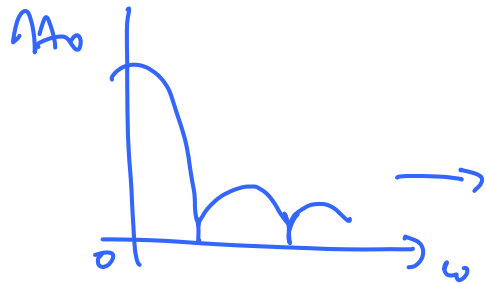
Exercise : Obtain the mag. freq response of $H_k(z)$ and plot them

HINT: Verify:

$$|H_0(z)| = \left| \frac{\sin(M\omega/2)}{\sin(\omega/2)} \right| \quad \text{--- (1)}$$

$$H_k(e^{j\omega}) = H_0(e^{j(\omega - 2\pi k/M)}) \quad \text{--- (2)}$$

Use (1) & (2) to sketch the freq. responses.



Time domain descriptions of multirate filters

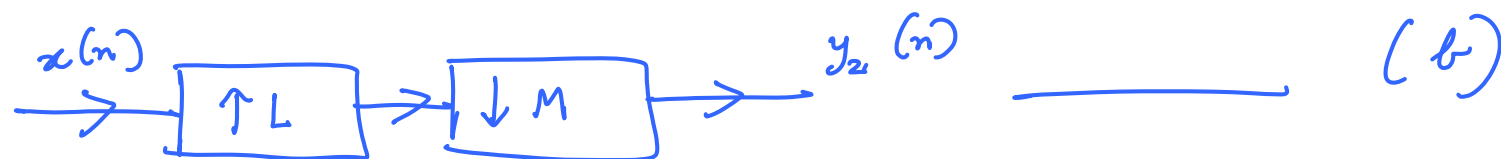
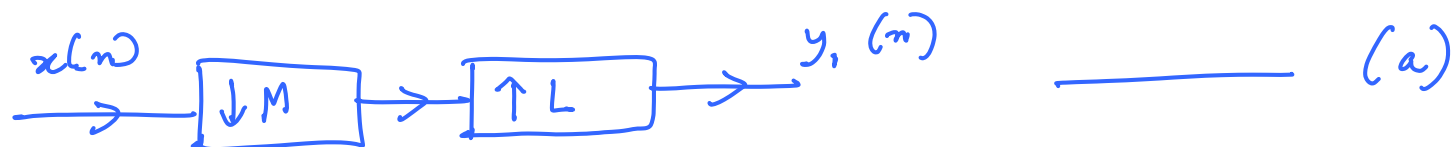
$$y(n) = \left\{ \begin{array}{l} \sum_{k=-\infty}^{\infty} x(k) h(nM - k) \\ \sum_{k=-\infty}^{\infty} x(k) h(n - kL) \\ \sum_{k=-\infty}^{\infty} x(k) h(nM - kL) \end{array} \right.$$

M-fold decimator

L-fold expander/inter

$\frac{M}{L}$ fold decimator

Interconnecting Systems



Question: Are (a) & (b) equivalent?

Theorem: The systems (a) & (b) are equivalent if L and M are relatively prime. i.e., $\gcd(L, M) = 1$

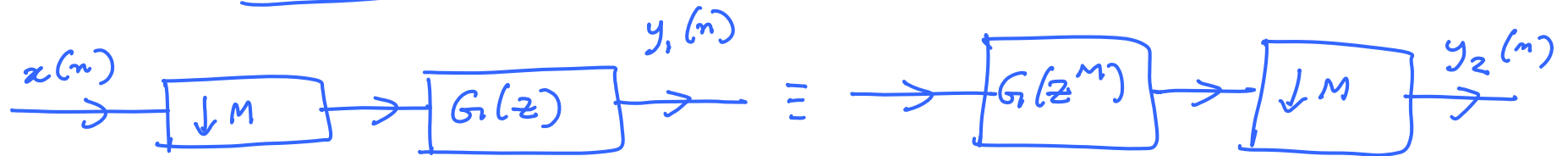
Sketch of proof:

$$Y_1(z) = \frac{1}{M} \sum_{k=0}^{M-1} X(z^{L/M} \underbrace{w_M^k}_{\text{red underline}})$$

$$Y_2(z) = \frac{1}{M} \sum_{k=0}^{M-1} X(z^{L/M} \underbrace{w_M^{kL}}_{\text{red underline}})$$

Claim: The set of numbers $S_1 = \{w_M^k\}_{0 \leq k \leq M-1}$ distinct M^{th} roots of unity is equal to the set of numbers $S_2 = \{w_M^{kL}\}_{0 \leq k \leq M-1}$ iff $\gcd(L, M) = 1$

Noble identities



Pathological Case : Suppose $G(z)$ is irrational i.e., $G(z) = z^{-1/2}$
 (fractional delay filter)

If $x(n)$ is such that $x(2n) = 0$ $M = 2$
 $y_1(n)$ is not zero for all n ; $y_2(n) = 0 \forall n$
 necessarily

Proof: (Non-irrational case)

$$\begin{aligned} \text{(a)} \quad Y_2(z) &= \frac{1}{M} \sum_{k=0}^{M-1} X\left(z^{\frac{1}{M}} \omega_M^k\right) G\left(\left(z^{\frac{1}{M}} \omega_M^k\right)^M\right) \\ &= \frac{1}{M} \sum_{k=0}^{M-1} X\left(z^{\frac{1}{M}} \omega_M^k\right) G(z) \quad \text{--- (1)} \end{aligned}$$

$$Y_1(z) = \frac{1}{M} \sum_{k=0}^{M-1} X\left(z^{\frac{1}{M}} \omega_M^k\right) G(z) \quad \text{--- (2)}$$

$$Y_1(z) = Y_2(z)$$

$$(r) \quad \text{III/IV} \quad Y_4(z) = G(z^L) X(z^L) \quad \text{—————} \quad (3)$$

$$Y_3(z) \stackrel{=}{=} X(z) G(z) \rightarrow \boxed{\uparrow L} \rightarrow$$
$$= X(z^L) G(z^L) \quad \text{—————} \quad (4)$$

$$Y_3(z) = Y_4(z)$$

Polyphase Representation

Inventor: Bellanger 1976

Idea: Create a set of lower order filters that work efficiently under multirate operations

Example

Basic Idea: $H(z) = \sum_{n=-\infty}^{\infty} h(n) z^{-n}$

Separate $H(z)$ into even and odd coeffs.

$$H(z) = \sum_{n=-\infty}^{\infty} h(2n) z^{-2n} + z^{-1} \sum_{n=-\infty}^{\infty} h(2n+1) z^{-2n}$$

$$\text{Let } E_0(z) = \sum_{n=-\infty}^{\infty} h(2n) z^{-n}$$

$$E_1(z) = \sum_{n=-\infty}^{\infty} h(2n+1) z^{-n}$$

$$H(z) = E_0(z^2) + z^{-1} E_1(z^2)$$

Example: Suppose (a) FIR

$$H(z) = 1 + 2z^{-1} + 3z^{-2} + 4z^{-3} \quad (\text{FIR case})$$

$$E_0(z) = 1 + 3z^{-1} + 4z^{-2}$$

$$E_1(z) = 2 + 4z^{-1}$$

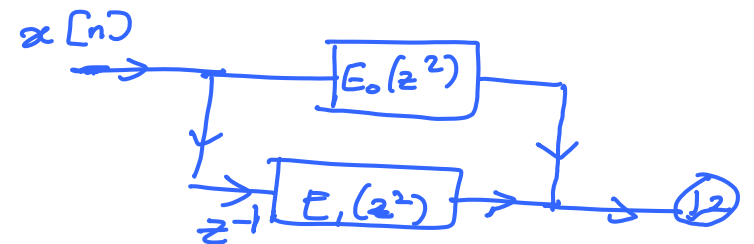
(4) 1.1.R Case

$$H(z) = \frac{1}{1 - \alpha z^{-1}} = \frac{1}{1 - \alpha^2 z^{-2}} + \frac{\alpha z^{-1}}{1 - \alpha^2 z^{-2}}$$

$$E_0(z) = \frac{1}{1 - \alpha^2 z^{-1}} \quad E_1(z) = \frac{\alpha}{1 - \alpha^2 z^{-1}}$$

$z^{-1} E_1(z^2)$

Having seen the basic trick, we would like to extend this idea for any integer M .



$$\begin{aligned}
 H(z) = & \sum_{n=-\infty}^{\infty} h(nM) z^{-nM} \\
 & + z^{-1} \sum_{n=-\infty}^{\infty} h(nM+1) z^{-nM} \\
 & \vdots \\
 & + z^{-(M-1)} \sum_{n=-\infty}^{\infty} h(nM+M-1) z^{-nM}
 \end{aligned}$$

This can be compactly written as

$$H(z) = \sum_{l=0}^{M-1} z^{-l} E_l(z^M) \quad \text{(Type 1 poly phase)}$$

$$E_l(z) = \sum_{n=-\infty}^{\infty} e_l(n) z^{-n}, \quad e_l(n) = h(nM+l)$$

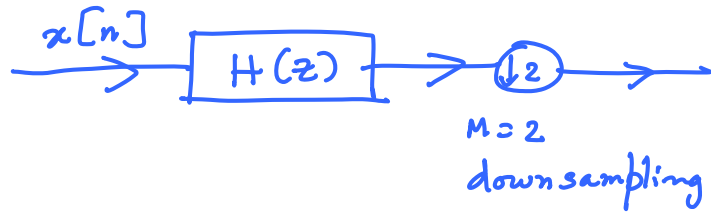
$0 \leq l \leq M-1$

Exercise : Show that

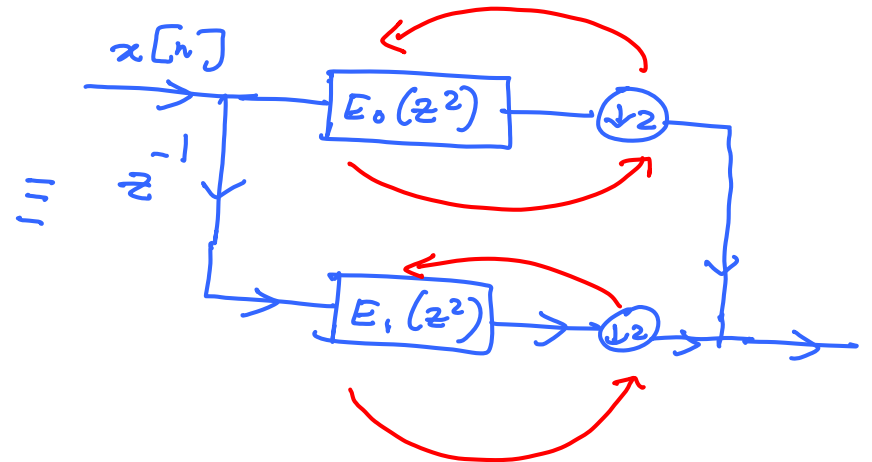
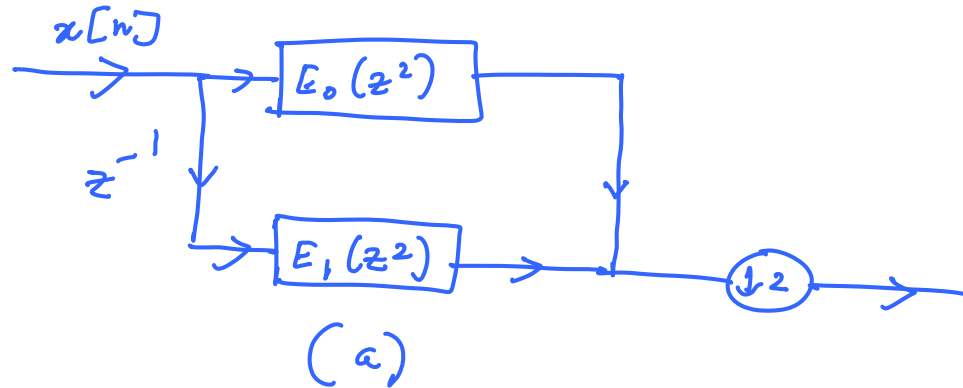
$$H(z) = \sum_{l=0}^{m-1} z^{-(m-1-l)} R_l(z^m) \quad (\text{Type 2 polyphase representation})$$

where $R_l(z^m) = E_{m-1-l}(z)$ (Permutations of $E_l(z)$)

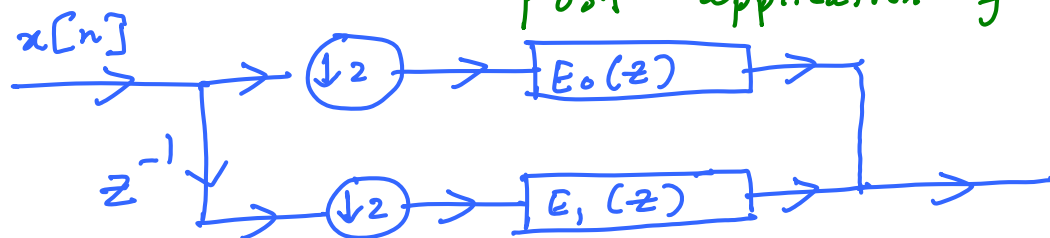
Efficient structures for decimation & interpolation filters



Let us do a Type 1 polyphase decomposition for $H(z)$



Post application of Noble Identities



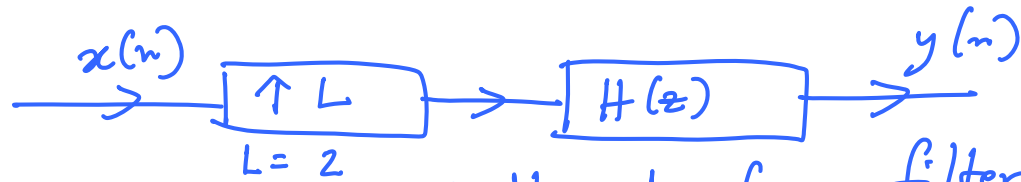
(b)

In the structure (a), because of down sampling, we consider only even numbered samples. During odd instants of time, the unit is just resting \Rightarrow inefficient resource utilization

Consider structure (b): Suppose n_0 and n_1 are the orders of $E_0(z)$ and $E_1(z)$ / $N+1 = n_0 + n_1 + 2$

The multipliers & adders in each of $E_l(z)$ $l = 0, 1$ have 2 units of time for doing their job & continuously operative. (No resting time)

Interpolation filters



Clearly a direct implement of a filter post upsampling is inefficient because 'at least' 50% of the time ($L=2$)

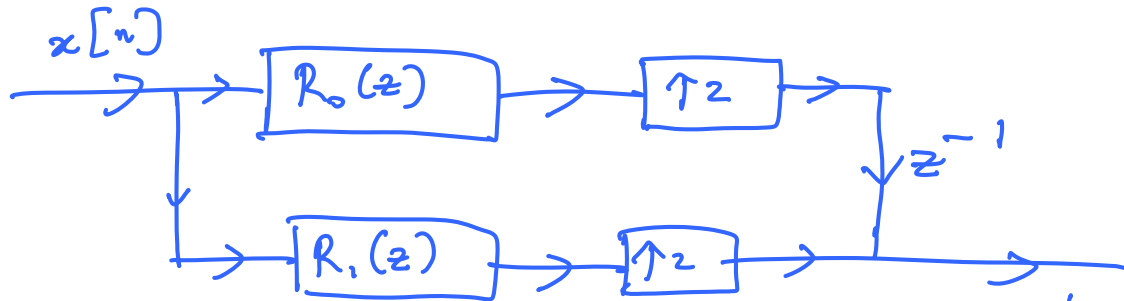
We are filtering zeros.

These multipliers "not resting" must complete the job in $\frac{1}{2}$ the time because o/p of the delay elements will change with time.

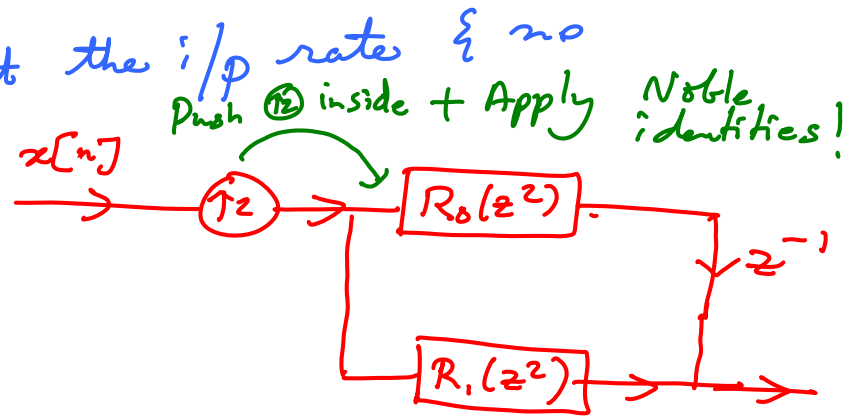
Using type 2 polyphase decomposition,

$$H(z) = R_1(z^2) + z^{-1} R_0(z^2)$$

$R_l(z)$ $l = 0, 1$ are operating at the i/p rate & no multiplier is resting
 Push $\uparrow 2$ inside + Apply Noble identities!



Efficient Structure for interpolation filters



Linear Phase FIR decimation filters

Let us suppose $H(z) = \sum_{n=0}^N h(n) z^{-n}$ such that

$$h(n) = h(N-n).$$

Let us investigate how symmetry in $h(n)$ reflects into the polyphase components.

Example: (a) Let $N = 4$

2 phase decomposition

$$E_0(z) = 1 + 4z^{-1} + z^{-2}$$

$$E_1(z) = 2 + 2z^{-1}$$

$$H(z) = 1 + 2z^{-1} + 4z^{-2} + 2z^{-3} + z^{-4}$$

Each of the polyphase component filters are symmetric!

(b) Suppose $N=5$

$$H(z) = 1 + 2z^{-1} + 4z^{-2} + 4z^{-3} + 2z^{-4} + z^{-5}$$

$$E_0(z) = 1 + 4z^{-1} + 2z^{-2}$$

$$E_1(z) = 2 + 4z^{-1} + z^{-2}$$

} $E_0(\cdot)$ & $E_1(\cdot)$ are
"mirrors" of each other!

Exercise: Devise an efficient architecture to exploit the mirror/symmetric properties of polyphase components in decimation & interpolation filters.